

Many-to-Many Shortest Paths

Using Highway Hierarchies

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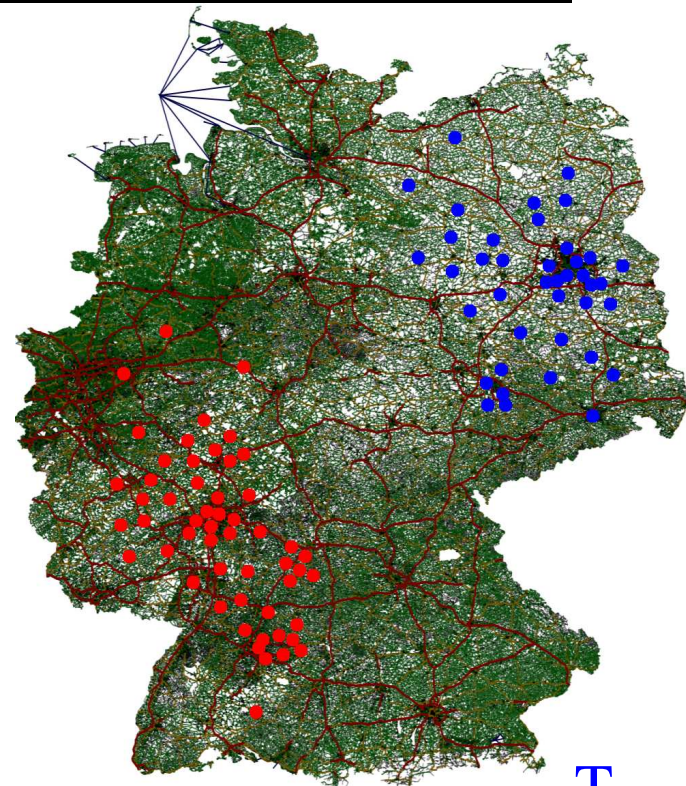
<http://algo2.iti.uka.de/schultes/hwy/>

New Orleans, January 6, 2007

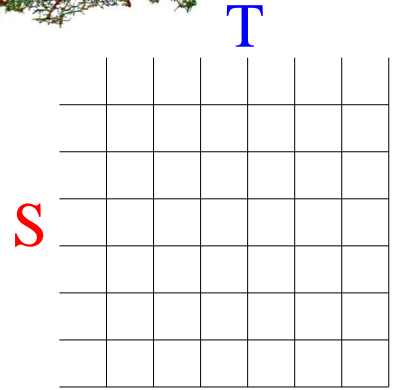
Many-to-Many Shortest Path Problem

Given:

- graph $G = (V, E)$
- set of source nodes $S \subseteq V$
- set of target nodes $T \subseteq V$



Task: compute $|S| \times |T|$ distance table containing the shortest path distances

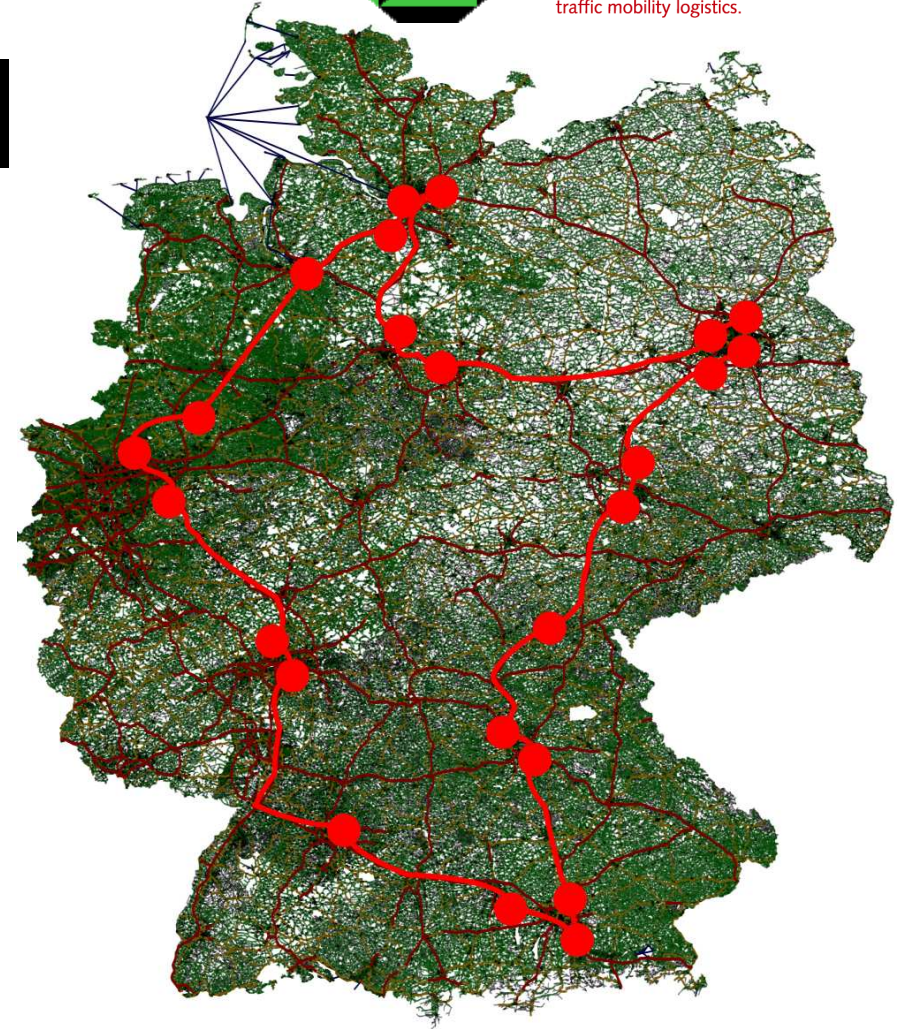


Here: concentrate on road networks

Applications

□ Logistics

- vehicle routing problem
- input for **traveling salesman** solver

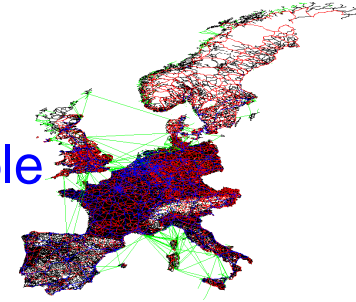


□ Preprocessing for Point-to-Point Techniques

- Precomputed Cluster Distances [\[MaueSandersMatijevic2006\]](#)
- **Transit Node Routing** [\[next talk\]](#)

Simple Solutions

Example: 10 000 × 10 000 table
in Western Europe

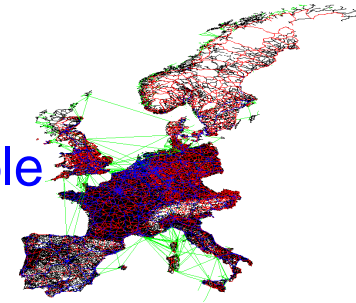


- apply SSSP algorithm $|S|$ times
(e.g. **DIJKSTRA**) $\approx 10\,000 \times 10\text{ s} \approx$ one day
- apply P2P algorithm $|S| \times |T|$ times
(e.g. **highway hierarchies**¹) $\approx 10\,000^2 \times 1\text{ ms} \approx$ one day

¹requires about 15 minutes preprocessing time

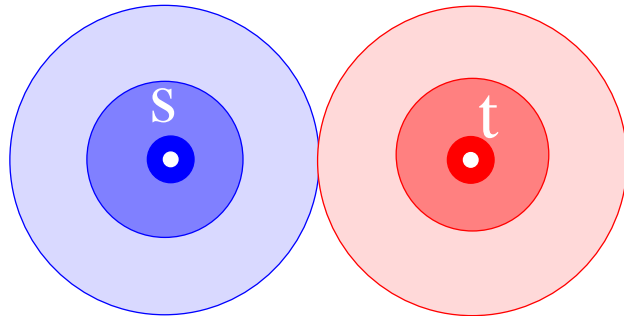
Our Solution

Example: 10 000 × 10 000 table
in Western Europe



- many-to-many algorithm
based on highway hierarchies¹

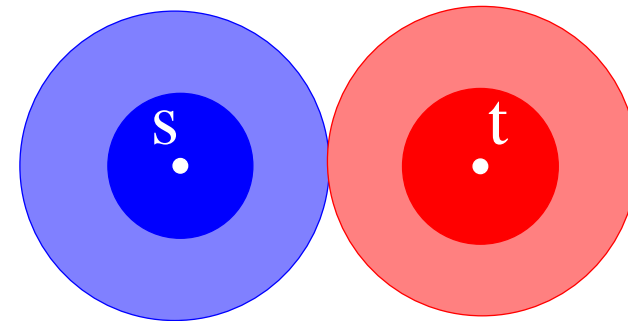
≈ one minute



¹requires about 15 minutes preprocessing time

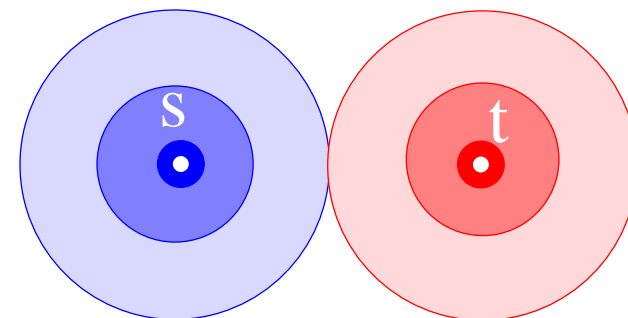
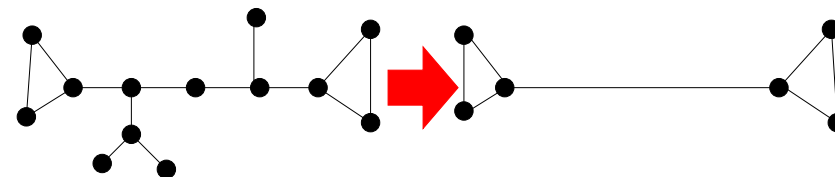
Highway Hierarchies²

- complete search within a local area
- search in a (thinner) highway network



= minimal graph that preserves all shortest paths

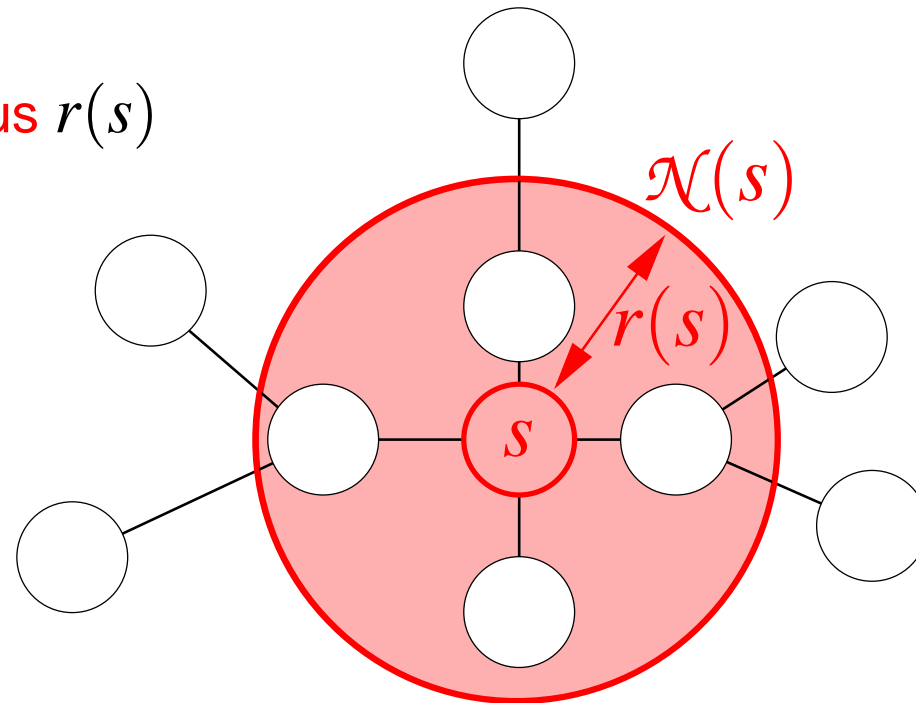
- contract network, e.g.,
- iterate \rightsquigarrow highway hierarchy



²presented at ESA 2005 and ESA 2006

Local Area

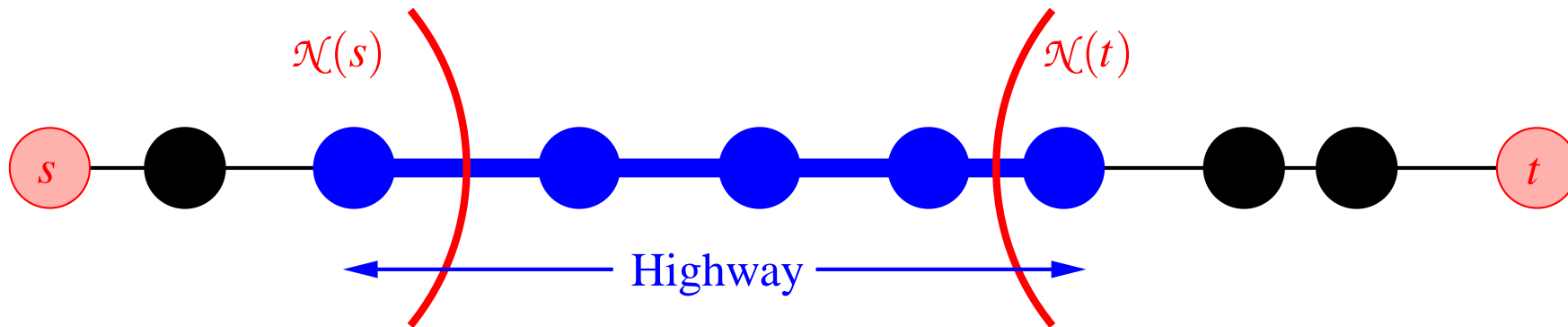
- choose **neighbourhood radius** $r(s)$
(by a heuristic)



- define **neighbourhood** of s

$$\mathcal{N}(s) := \{v \in V \mid d(s, v) \leq r(s)\}$$

Highway Network



Edge (u, v) belongs to **highway network** *iff* there are nodes s and t s.t.

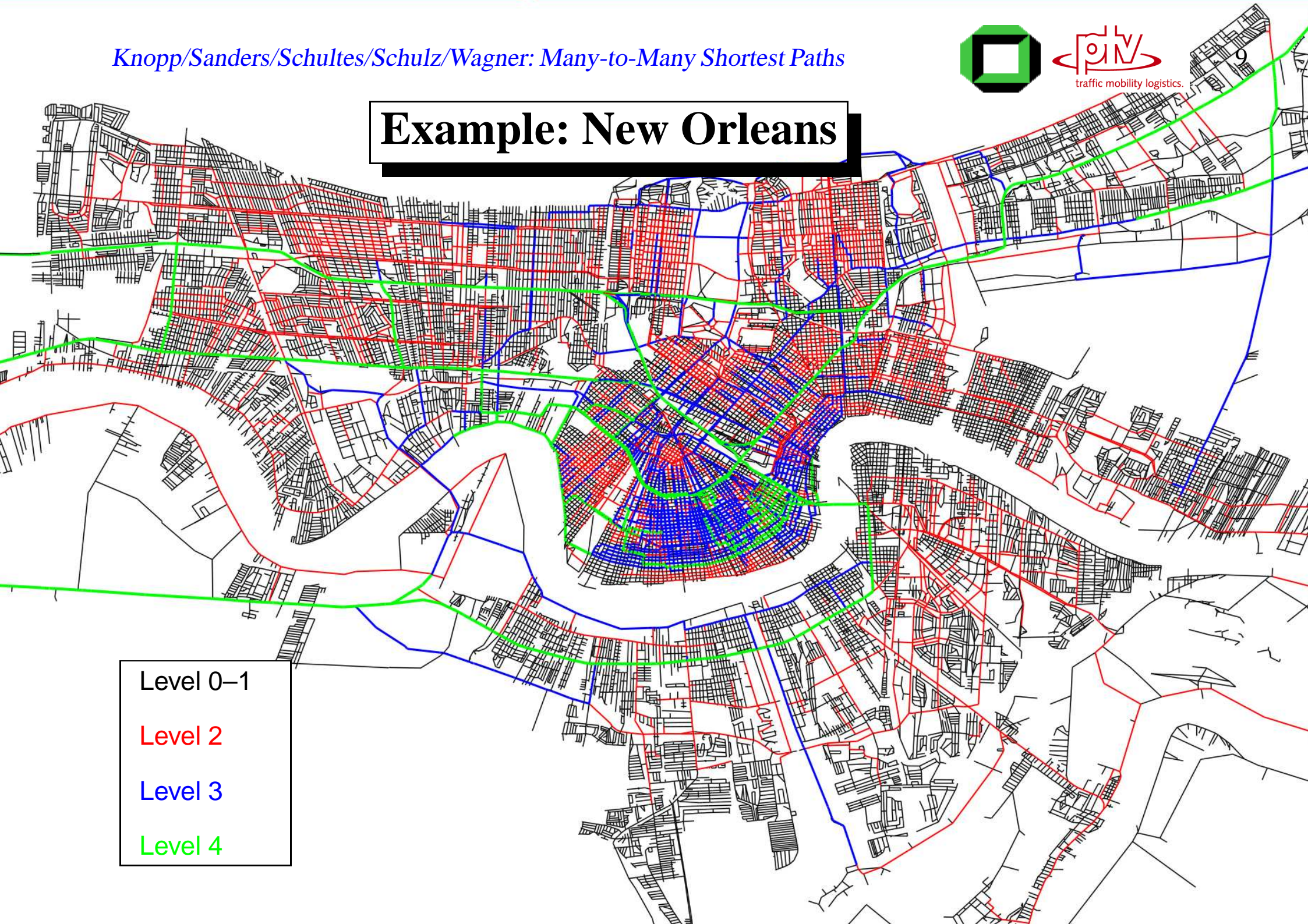
(u, v) is on the shortest path from s to t

and

(u, v) is not entirely within $\mathcal{N}(s)$ or $\mathcal{N}(t)$



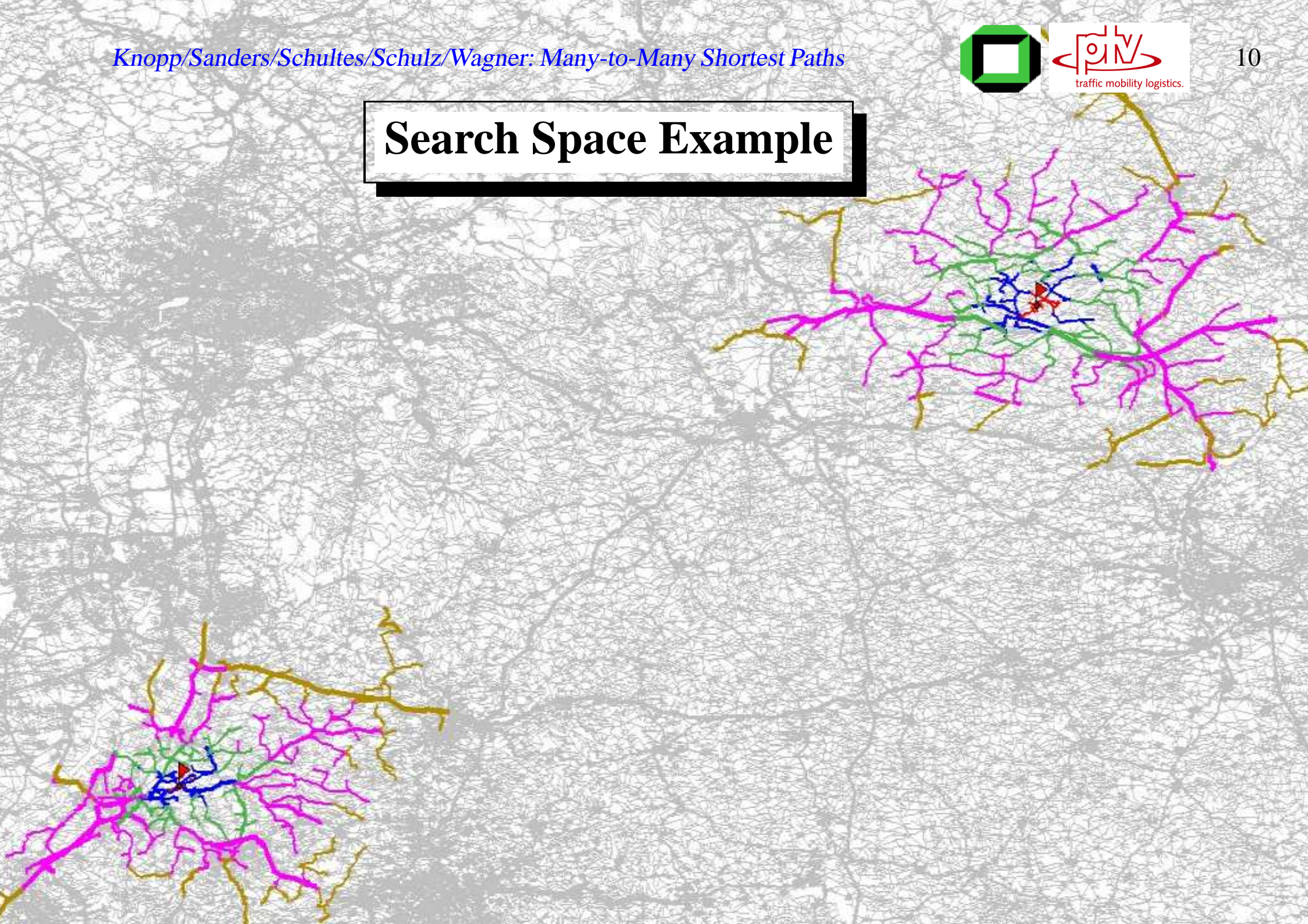
Example: New Orleans



- Level 0-1
- Level 2
- Level 3
- Level 4



Search Space Example

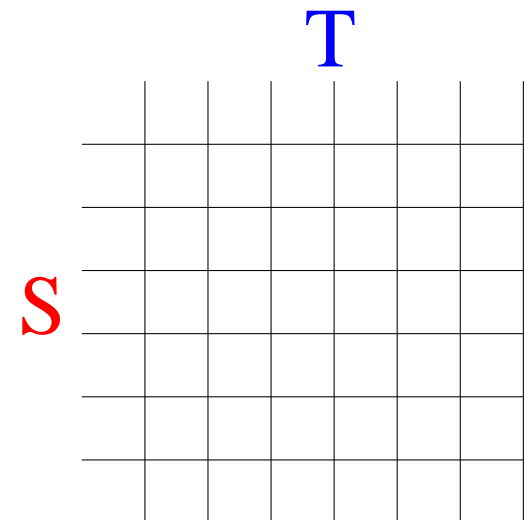


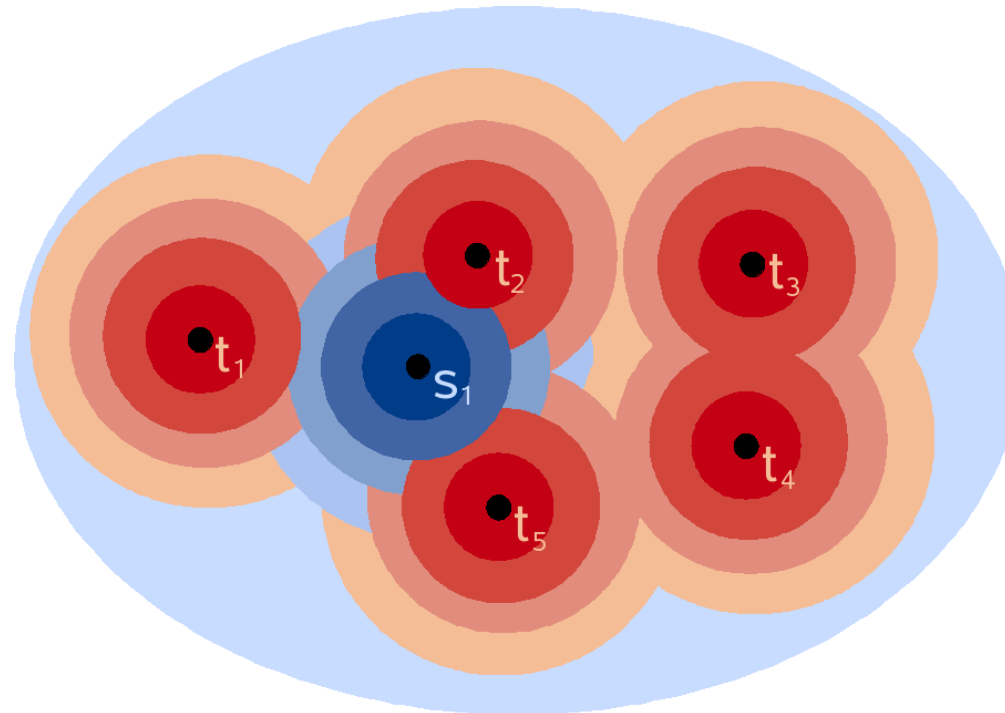
Main Idea

- instead of $|S| \times |T|$ **bidirectional** highway queries
- perform $|S| + |T|$ **unidirectional** highway queries

Algorithm

- maintain an $|S| \times |T|$ table D of **tentative distances**
(initialize all entries to ∞)



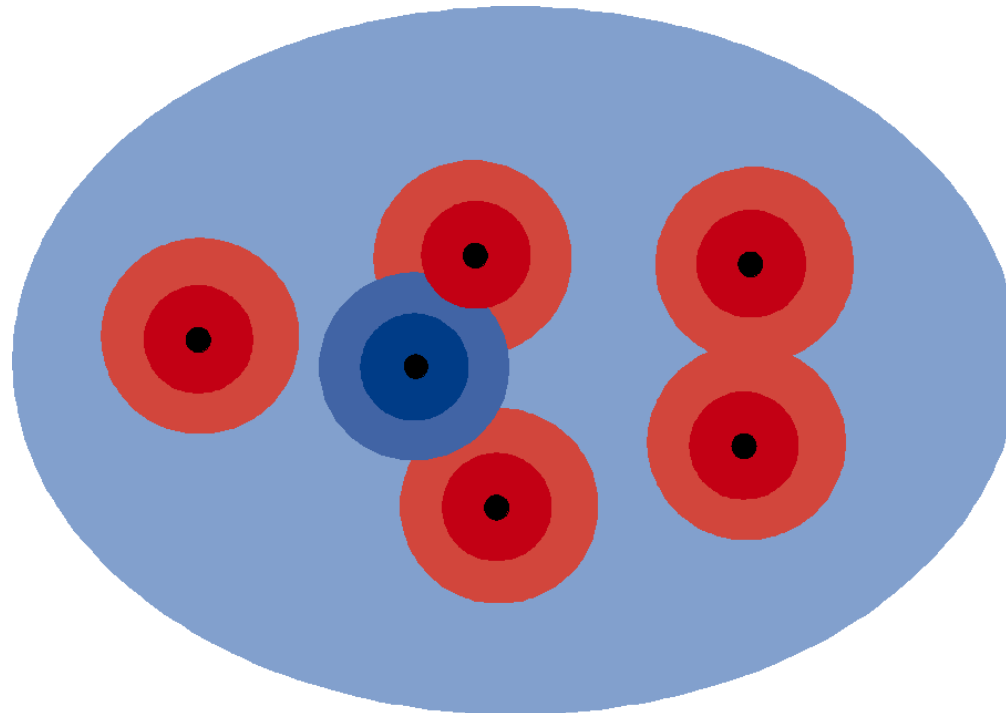


- for each $t \in T$, perform **backward search** up to the top level,
store search space entries $(t, u, d(u, t))$
- arrange search spaces: create a bucket for each u
- for each $s \in S$, perform **forward search** up to and **including** the top level,
at each node u , **scan all entries** $(t, u, d(u, t))$ and
compute $d(s, u) + d(u, t)$, update $D[s, t]$

Asymmetry

for large distance tables, most time spent on **bucket scanning**

Solution: use **less levels** \rightsquigarrow strengthen the **asymmetry**

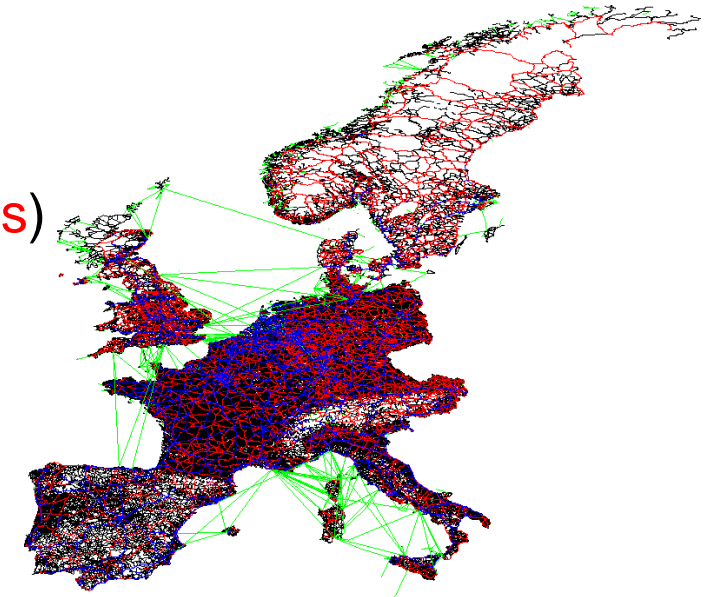


- backward search spaces get smaller \rightsquigarrow **less bucket entries**
- forward search spaces get bigger

Experiments

Input:

- Western European road network (**18 million nodes**)
- random source/target node sets



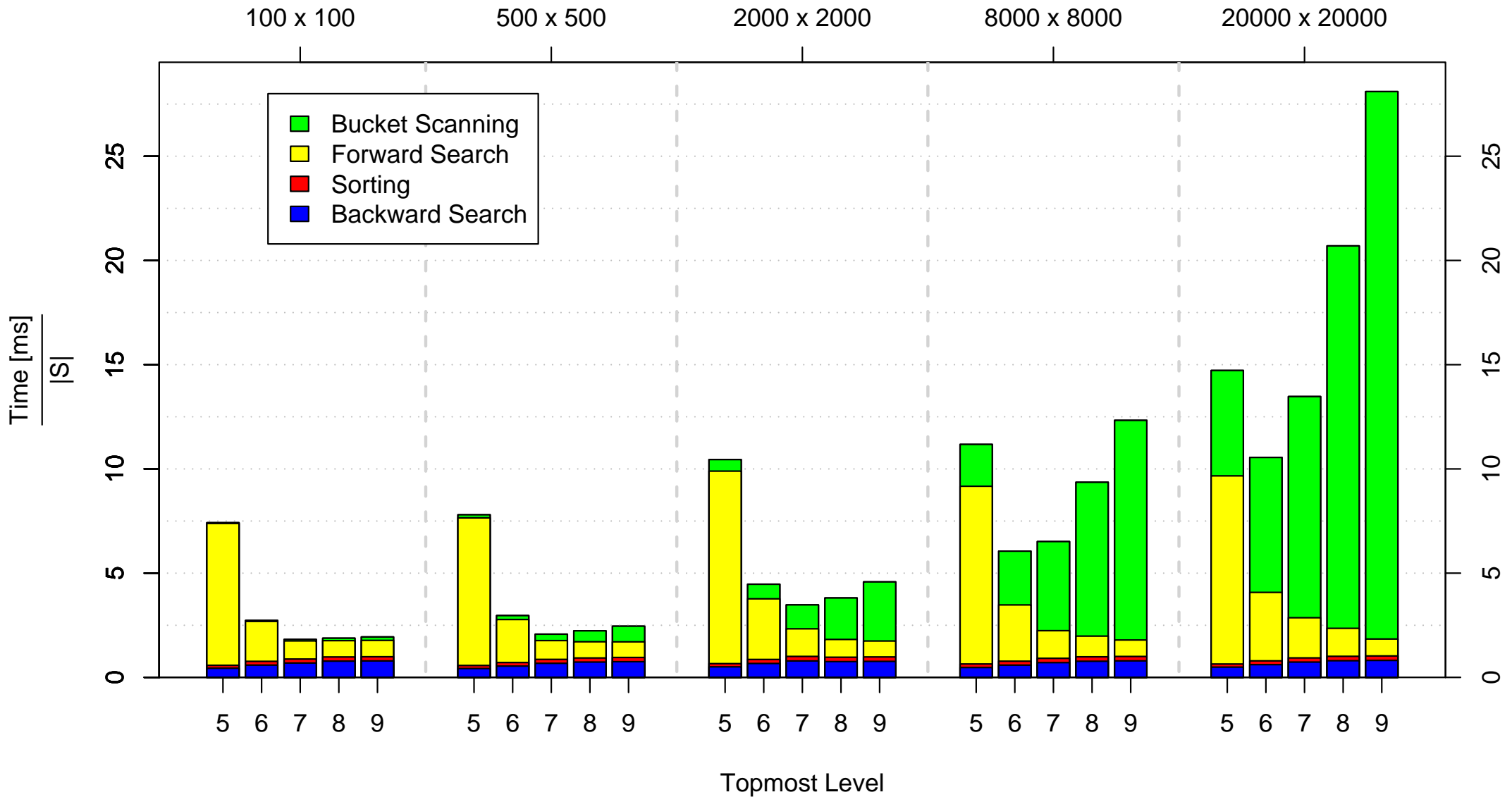
Results:

table size	time	speedup (\leftrightarrow DIJKSTRA)
1 000 \times 1 000	2.5 s	4 680
10 000 \times 10 000	58 s	2 017

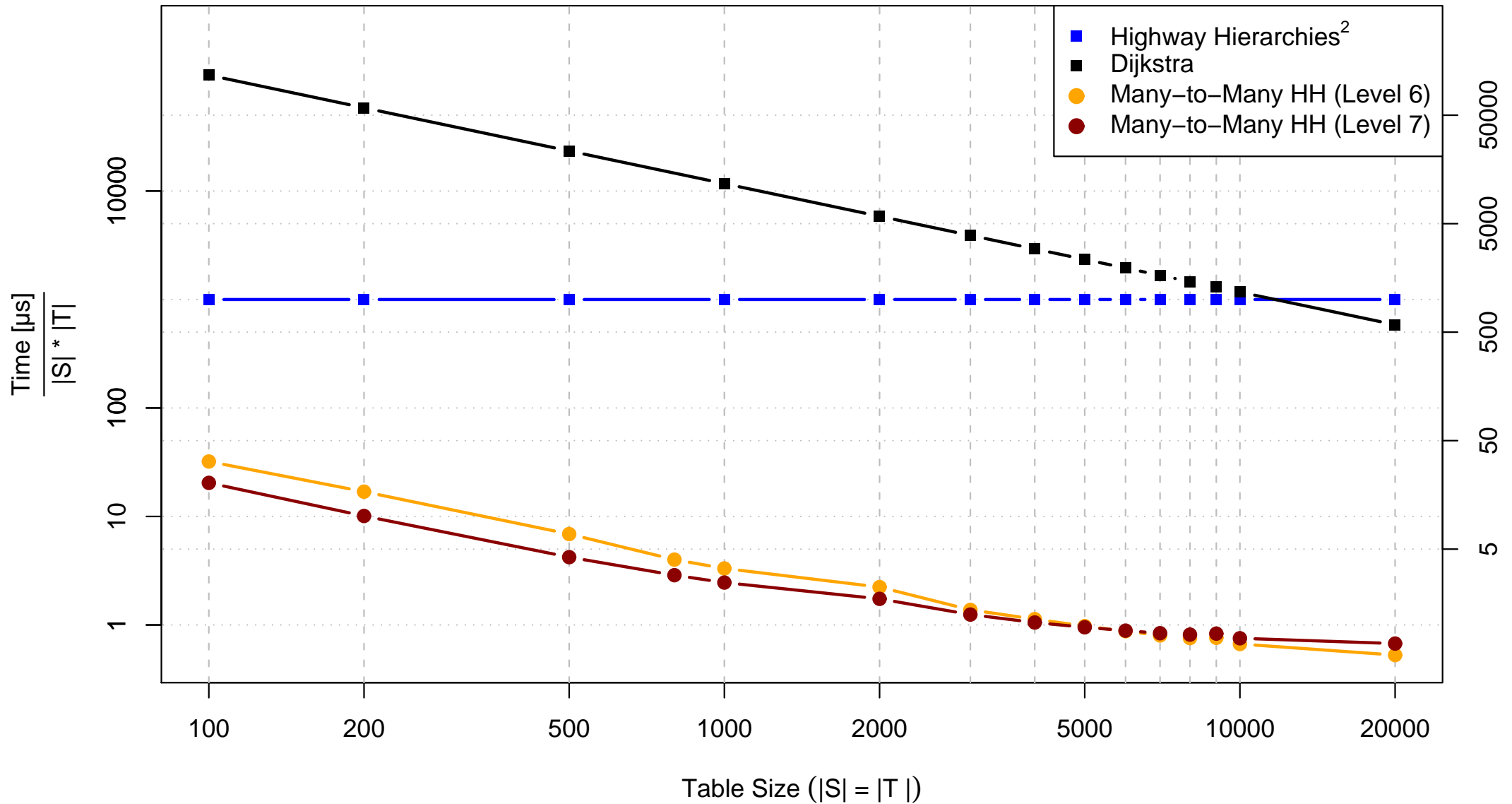
Break Even Point (w.r.t. preprocessing costs): **table size 100 \times 100**

Real-World Instances: **similar performance**

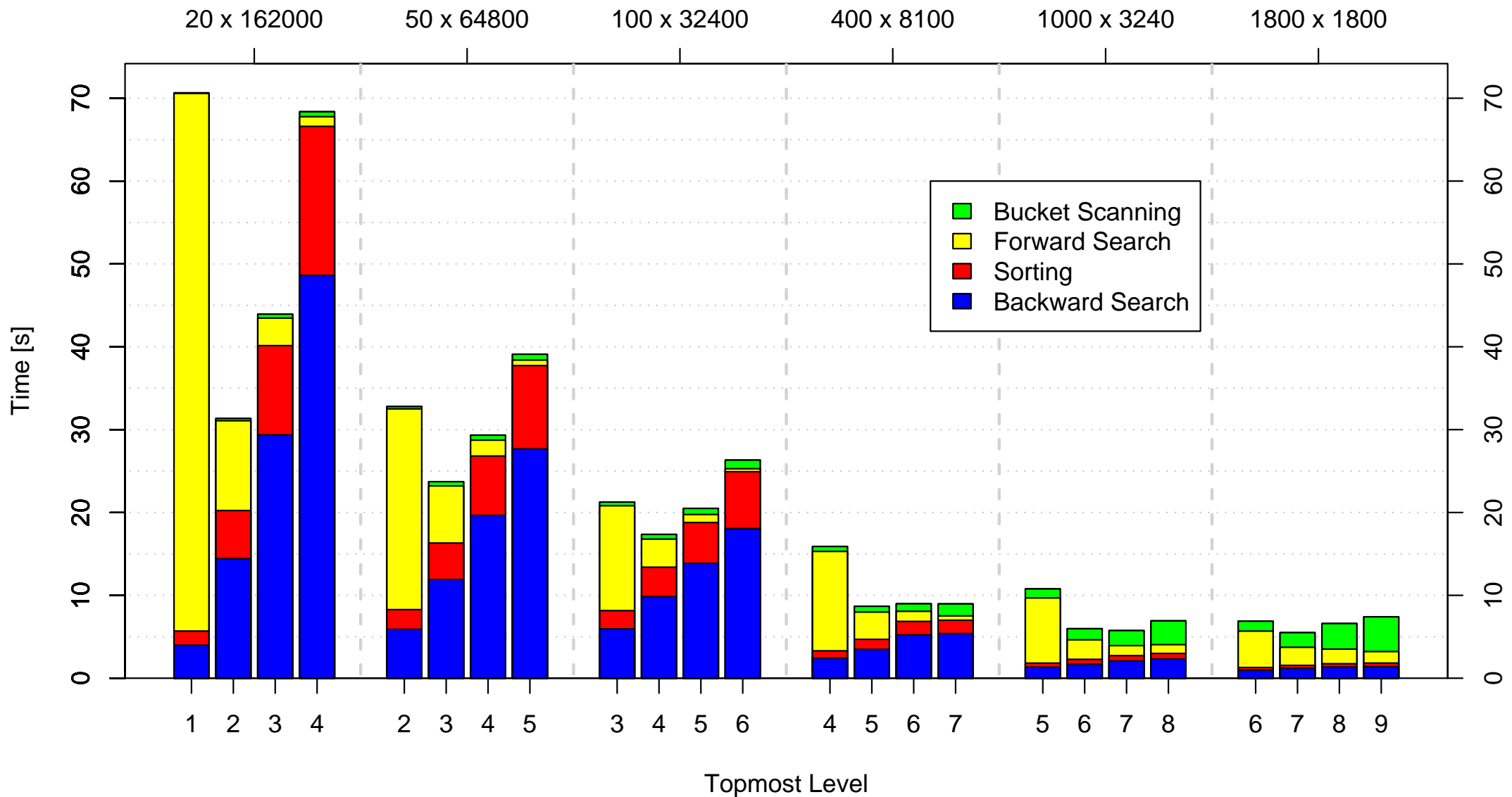
Symmetric Instances



Comparisons

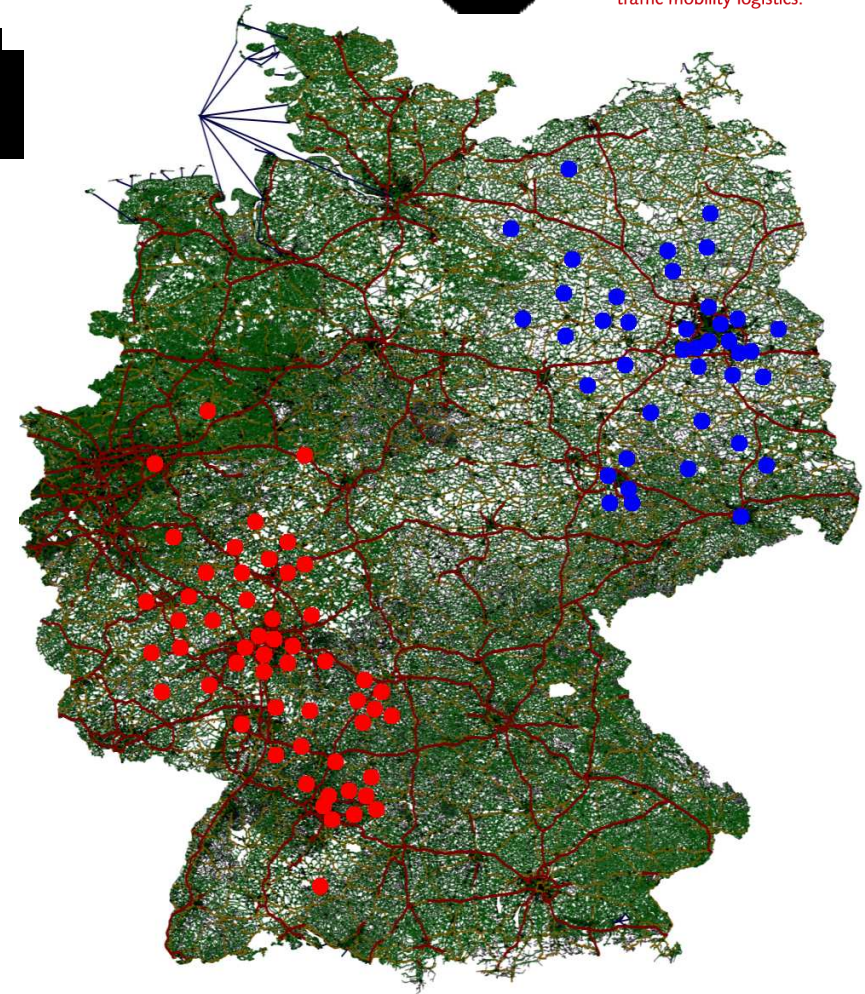


Asymmetric Instances



Summary

□ very efficient solution to the many-to-many shortest path problem



□ requires little preprocessing time

≈ 15 minutes

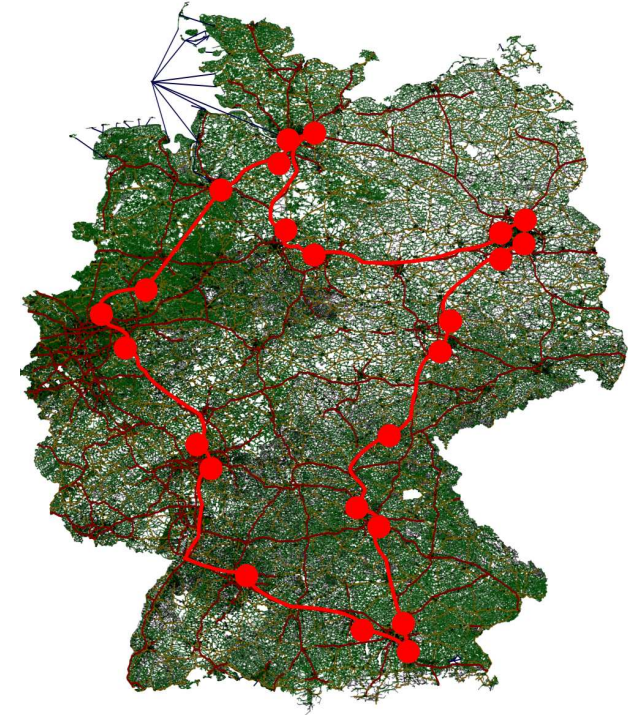
□ computes 10 000 × 10 000 table in

≈ 1 minute

(0.6 μs per entry)

Additional Issues

- outputting paths
- incremental computation
- parallelization



Future Work

- adapt** preprocessing to **specific** source/target node sets
- approach can be **generalized** to other

- **non-goal-directed**
- **bidirectional**

speedup techniques

